

INTRODUCTION

The TINI[®] platform¹ JVM is powerful enough for very high-level tasks such as serving web pages or remote data logging. However, the Java language only provides limited low-level access necessary for hardware communication. For speed and determinism sake, tasks of this nature are better implemented at an operating system level.

The TINI I/O Subsystem enables communication with devices that require faster communication than the JVM allows. It provides a mechanism for assembly language drivers to act as an intermediary between hardware and the JVM in a standard way. Additional drivers can be installed for communication with new devices. This article examines the TINIOS I/O subsystem, explains how to write a driver as a native library, and provides a simple example in the form of a pipe driver.

GETTING STARTED

Before we begin, you need the TINI SDK version 1.02e or higher. Versions prior to 1.02e did not publicly expose the `com.dalsemi.comm.NativeComm` class that is necessary for driver communication.

This document assumes some familiarity with the TINI Native Interface (TNI) as well as familiarity with 8051 assembly, the `a390` assembler, and the `macro` preprocessor².

DRIVER BASICS

A driver in TINIOS is made up of two parts: a set of standard functions to link with TINIOS (the driver interface) and an interrupt or polled routine (the driver) to communicate with the device. Once the interface is registered with TINIOS, it can be accessed using the `com.dalsemi.comm.NativeComm` class.

Accessing a Driver in Java

A device driver is loaded in as a native library. Once loaded, the following `com.dalsemi.comm.NativeComm` methods can be used.

- `public static int open(int port, int stream);`
Opens a connection to the driver. `port` represents the driver's port number, which is chosen by the driver when it registers itself. `stream` can be set to `NativeComm.STREAM_STDIN` or `NativeComm.STREAM_STDOUT` (they cannot be logically OR-ed together), depending on if the handle represents an input or output stream. An integer handle is returned, which represents the stream in all other calls. The handle is actually a one-byte unique-identifier value from the I/O subsystem, but promoted to an integer for Java.
- `public static int close(int handle);`
Closes the stream associated with `handle`. Returns a 1 for success, or a 0 for failure.

¹ Information available at <http://www.ibutton.com/TINI>.

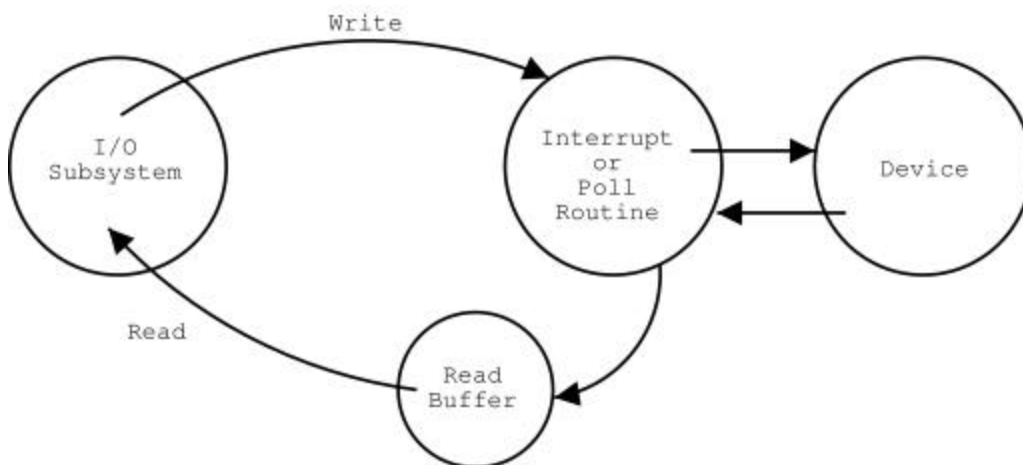
² The TINI SDK, including all tools, are available at <ftp://ftp.dalsemi.com/pub/tini>.

- `public static int read(int handle, byte [] arr, int timeout, boolean suspend);`
- `public static int read(int handle, byte [] arr, int length, int offset, int timeout, boolean suspend);`
Reads from the opened stream. `handle` must represent a stream opened for reading. `arr` is the byte array buffer for the input. `length` is the size of the input buffer. `offset` represents the first element of `arr` to be used for the read. `timeout` is the time in milliseconds allowed for the read to complete before it returns. `suspend` should be set to true to ignore the timeout and suspend until the read completes, or false to use the timeout. Function returns the number of bytes read.
- `public static void write(int handle, byte [] arr);`
- `public static void write(int handle, byte [] arr, int offset, int length);`
Writes to an opened output stream. `handle` must represent a stream opened for writing. `arr` is the output array. `offset` is the first element of `arr` to be written. `length` is the number of bytes to write.
- `public static int ioctl(int handle, byte [] arr, int offset, int length, int timeout);`
Sends an `ioctl` call to the driver. `handle` is any valid stream handle. `arr` is an array of parameters. `length` is the number of parameters. All other arguments should be ignored.
- `public static int available(int handle);`
Returns the number of bytes available without blocking. `handle` represents an input stream.

DRIVER PHILOSOPHY

Figure 1 shows the design of a TINIOS device driver.

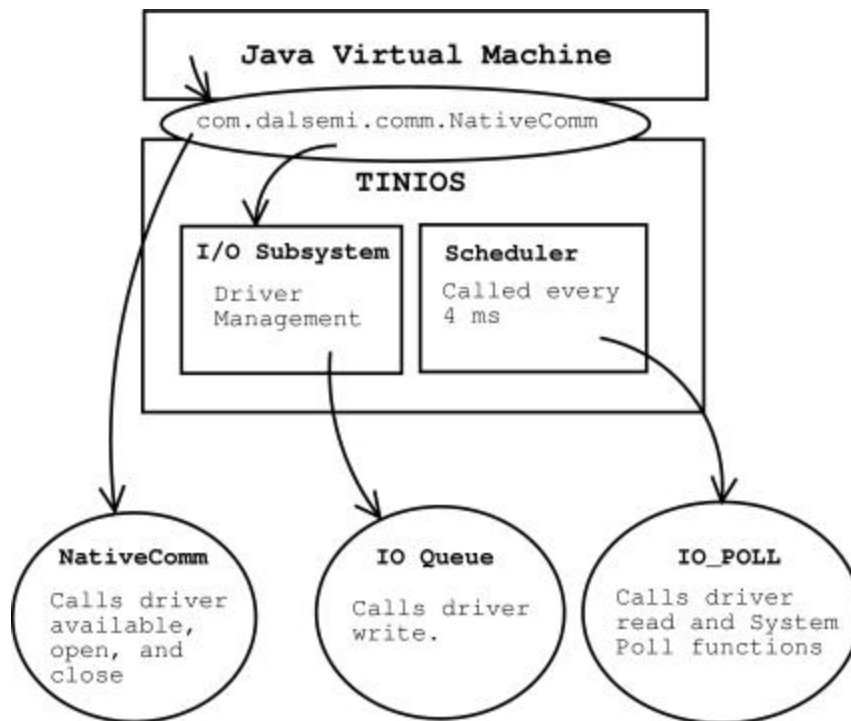
Figure 1. I/O DRIVER PHILOSOPHY



The driver can store incoming data in some sort of buffer in the event a read cannot be handled immediately. A write is directly fed from the driver interface to the driver.

The driver interface consists of implementations of `open`, `close`, `read`, `write`, `available`, and `ioctl` functions. These can be grouped together by what portion of the I/O subsystem invokes them, as shown by Figure 2.

Figure 2. TINIOS I/O DESIGN



NativeComm

The `com.dalsemi.comm.NativeComm` class binds the Java Virtual Machine and the I/O subsystem. It contains static Java methods for calling the driver interface functions. While the `read` and `write` driver calls are queued by the I/O subsystem, the `open`, `close`, `available`, and `ioctl` are invoked directly from the `NativeComm` methods. Program execution continues uninterrupted after these functions complete.

- `open`: Called when the user tries to open a device for input or output. Driver and device initialization should be performed in this function. Returns a success/error value.
- `close`: Called when the user closes a device. Driver and device shutdown procedures should be performed in this function. Returns a success/error value.
- `available`: Called to check the amount of available data that can be read from the driver. Returns an integer count.
- `ioctl`: Called for functionality not directly supported by the driver interface. Returns a success/error value.

I/O Queue

The I/O subsystem handles writes differently than the above methods.

- 1) The I/O subsystem receives the `write` call, checks if the handle is open for writing, and queues it if the device is busy.
- 2) The driver `write` function is passed a buffer and a byte count to write. The driver `write` is a catalyst—it does not perform the write operation, but instead prepares the driver for writing and returns.
- 3) The I/O subsystem puts the executing Java thread to sleep.
- 4) The driver performs the hardware write. When it is finished, it calls the `IO_WriteFree` function to inform the I/O Subsystem it has completed.
- 5) The Java thread is awoken, and execution resumes.

It is important to note that the driver `write` does not write to the hardware, but only prepares the write for asynchronous execution by the driver. The safest approach is to have a poll routine check for the completion of the write, calling `IO_WriteFree` when it detects completion. If `IO_WriteFree` is called from the `write` function, it attempts to wake the Java Thread before it is put to sleep, and the **thread suspends indefinitely**. Also, TINI does not have any way of protecting operating system state from being modified inside of an interrupt service routine. Therefore, `IO_WriteFree` should **not** be called from an ISR.

IO_POLL

The I/O subsystem must check on registered drivers continuously. To perform these checks, the `IO_POLL` function is called by the scheduler every 4ms³. `IO_POLL` manages the status of pending reads and writes. It is possible to register your own callback functions to be called on every `IO_POLL`. These do not necessarily need to be I/O related and can therefore serve any purpose.

The `read` driver interface function is called by `IO_POLL` as well. A driver `read` behaves very differently than a driver `write`.

- 1) The read is queued by the I/O subsystem.
- 2) The Java thread is put to sleep.
- 3) `IO_POLL` calls the driver `read` function. The driver `read` uses an intermediate buffer of its own to return up to 250 bytes back to the I/O subsystem. The driver `read` copies the amount it has available into the buffer and exits.
- 4) Step 3 repeats until the amount requested is read in, or a timeout occurs in `IO_POLL` (if a timeout is enabled).
- 5) The Java thread is awoken by the I/O subsystem upon completion of the read or a timeout.

Note the differences between write and read. While a driver `write` is called once to initiate a write to a device, a driver `read` is called repeatedly by the I/O subsystem to poll the data. Also, while a write operation requires an `IO_WriteFree` for termination, a read operation is terminated automatically by the I/O subsystem on completion.

³ More or less 4ms. TINIOS is not real-time.

IMPLEMENTATION DETAILS

Driver Portion

If the device uses an interrupt, use `System_Install_Interrupt` to install your interrupt service routine, or `System_Register_External_Interrupt` to chain to the external interrupt line. If the device requires a poll-based routine, use `System_RegisterPoll` to have your routine called every 4ms by `IO_POLL`.

Device writes should be handled within the driver, calling `IO_WriteFree` on completion. It is important to note that `IO_WriteFree` cannot be called by the driver write interface function or by any interrupt service routine.

Driver Interface Registration

To add a driver to the TINIOS, you need to register your driver with the I/O subsystem using the function `System_Register_Driver`. A driver number between `IOSYS_USER_DRIVER` and `IOSYS_MAX_DRIVER` must be assigned to the driver⁴. This number is used when opening and closing the port.

It is recommended that you perform this step in the initialization function of your native library. The driver will not be available until this step is complete.

Open

The driver open function is called when a thread attempts to open a driver. One of the microcontroller registers (`R2_B0`) is set to the handle assigned to the port. Any state and device initialization should be performed here.

Clear the accumulator to denote success, or an exception number (defined in `apiequ.inc`) to throw an exception. All other registers must be restored to their prior state before returning.

Close

The driver close function is called when a thread closes the driver. Resident state cleanup and device shutdown should be handled here.

Clear the accumulator to denote success, or an exception number (defined in `apiequ.inc`) to throw an exception. All other registers are restored to their prior state before returning.

Read

Once a device read is initiated, the driver read function is called repeatedly. It is not passed any arguments, and it returns a pointer to a structure, described by Figure 3.

Figure 3. READ FUNCTION RETURN VALUE



⁴ TINIOS supports 16 I/O drivers. Drivers 0 to 7 are allocated to the system, and 8 to 15 are allocated to the user.

The first byte is the size of the buffer being returned, with a maximum of 250 bytes. The next three bytes are reserved, and should be set to 0. The remaining bytes are the data from the resident portion's read buffer. *Do not create a new buffer on every read call!* Instead, create a buffer in the `open` function and free it in `close`. If speed is not an issue, you can use ephemeral state blocks (with `NatLib_GetEphemeralStateBlock` and `NatLib_RemoveEphemeralStateBlock`) to hold your state information. Otherwise, use indirects (acquired from `System_AcquireIndirectSemaphore`) to store the pointer to your state.

Clear the accumulator to denote success, or an exception number (defined in `apiequ.inc`) to throw an exception. All registers except the first data pointer must be restored before returning.

Write

The driver write is passed the application id, the thread id, the data length, and a pointer to the data to be written. All of this information must be passed to the resident portion and the `write` must be initiated.

Clear the accumulator to denote success, or an exception number (defined in `apiequ.inc`) to throw an exception. All other registers must be restored on completion.

Available

Available is called to query the amount of data that can be read from the driver without blocking. The value can be returned in R3:R2:R1:R0 of register bank 1. Clear the accumulator to denote success, or an exception number (defined in `apiequ.inc`) to throw an exception. All other registers are restored to their prior state before returning.

ioctl

The `ioctl` (input output control) function is used to make device specific calls. The first data pointer point to an argument array, whose length is stored in R5:R4.

Clear the accumulator to denote success, or an exception number (defined in `apiequ.inc`) to throw an exception. A return value can be put in R3:R2:R1:R0 of register bank 1. All other registers must be restored on completion.

Process Destruction

A driver can register a process destroy callback, using `System_RegisterProcessDestroyFunction`, to perform any closing and shutdown operation when a process exits. The process destroy callback is passed the identifier of the dying process in R0. Clear the accumulator to denote success, or an exception number to throw an exception. All other registers must be restored on completion.

Example: System Pipe

Provided in the appendices is a sample driver that implements a very simple pipe driver. The pipe driver allows for output from one process or thread to be read from another process or thread. It uses an Ephemeral State Block (ESB) to hold the circular buffer and the read function return buffer. One byte is used for the size count, and another is used for the offset of the first byte into the queue. Using a 256-byte buffer makes handling the rollover simple, as it occurs automatically when the appropriate registers overflow.

On startup, the native library initialization routine uses the process id from `System_GetCurrentProcessId` (which is assigned a value of 0 to 7) to set a bit in the reference counter value of the ESB. It then registers

a process destroy callback, which clears the bits of the reference counter. When the counter is 0, the driver performs cleanup operations and unregisters itself. This allows communication across process boundaries.

Memory past the circular buffer is used for the read return value. There is no need to worry about serializability of the return buffer, because only one driver `read` can be called at a time. The read function insures it copies less than 250 bytes before performing the copy.

The write function installs a poll routine in `IO_POLL` to write the data into the buffer as space becomes available. The poll routine performs the copy, storing intermediate state near the end of the ESB. The poll routine removes itself when the write completes.

Classes in the `com.dalsemi.pipeio` package implement a `java.io.InputStream` and `java.io.OutputStream` on top of the driver. This allows for java programs to access the driver as if it was another standard Java input or output stream.

CONCLUSION

While most embedded tasks can be handled by the TINI JVM, there are instances that require more greater and tighter constraints. These can be handled by adding native drivers to TINIOS. A native driver consists of a polled or interrupt handler, and a set of interface functions. They can be loaded inside of a native library, and TINI provides a standard interface to communicate with them. The main advantage of using a native driver is that most of the operating system tasks, like queuing multiple writes or blocking a process on a read, are offloaded onto TINIOS. This frees the developer to focus on device communication and not operating system semantics.

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Appendix A: driver.a51

```

#include(ds80c390.inc)
#include(apiequ.inc)
#include(tini.inc)
#include(tinimacro.inc)
#include(driver.inc)

;
; This driver implements a 255 byte circular buffer for
; communication across threads or processes. It user
; an ephemeral state block to hold the circular buffer
; and the return value for the read. (The structure
; is described in driver.inc).
;

;*****
;*
;* Function Name: Init_Pipe
;*
;* Description: Native Library initialization routine.
;*
;* Input(s): None.
;*
;* Outputs(s): a - 0 on success, non-zero for init exception.
;*
;* Notes: I am creating my state info in this function.
;*
;*
;*****

Init_Pipe:
;
; This function performs two operations. The first is
; to register the driver with the I/O subsystem. The second
; is to create state info for the pipe (the circular buffer,
; the read return value buffer, and the write state
; buffer).
;

Init_GetEphemeralStateBlock:
;
; Now create the state info. First, check if the
; Ephemeral State Block (ESB) exists.
;
mov   dptr, #QueueID           ; see if we have an ESB installed
lcall NatLib_GetEphemeralStateBlock

;
; Cache pointer away if it gets destroyed
;
mov   R5,dpl
mov   R6,dph
mov   R7,dpx

jnz   Init_MallocEphemeral    ;

```



```

;
; an ESB exists - some process
; has already called load since
; boot time
;
Init_StoreRetrieveEphemeral:
    ljmp    Init_IncrementReferenceCount

;
; No ESB exists! Dios Mio! Time to
; init all the data!
;

Init_MallocEphemeral:
;
; Install the process destroy function
;
mov     dptr, #Pipe_ProcessExit
lcall   System_RegisterProcessDestroyFunction

mov     R2, #LOW(PIPE_END)        ; Allocate 1024 bytes
mov     R3, #HIGH(PIPE_END)      ;
lcall   mm_Malloc                ;
jnz     Init_Done                ;

;
; The memory is cleared by mm_malloc, so there isn't
; any more initialization to do
;
mov     dpl1, dpl                ; copy address to
mov     dph1, dph                ; second dptr
mov     dpx1, dpx                ; handle is already in R3:R2

;
; Cache pointer away if it gets destroyed
;
mov     R5, dpl
mov     R6, dph
mov     R7, dpx

mov     dptr, #QueueID           ; first dptr is our identifier
lcall   NatLib_InstallEphemeralStateBlock

mov     dpl, R5
mov     dph, R6
mov     dpx, R7

;
; We need to now assign our driver a
; driver number and store it away
;
mov     a, #LOW(PIPE_DRVNUM)
mov     b, #HIGH(PIPE_DRVNUM)
lcall   add_dp1r1_16
lcall   System_IO_NextAvailableDriverNum
PUTX
mov     dpl, R5
mov     dph, R6
mov     dpx, R7

;

```

```

; Register the driver with the I/O subsystem.
; acc holds the driver num
;
lcall  Pipe_Init
push   acc
lcall  System_Register_Driver
;
; Disable non blocking write I/O
;
pop     acc
mov     b,#01
lcall  System_IO_EnableNonBlockingWrites

mov     dpl,R5
mov     dph,R6
mov     dpx,R7

```

Init_IncrementReferenceCount:

```

;
; mm_malloc should initialize the memory to zero.
; but we need to increment the reference count
;
mov     a,#LOW(PIPE_REFCOUNT)
mov     b,#HIGH(PIPE_REFCOUNT)
lcall  add_dpstr1_16
;
; Set the bit denoting the process id
;
lcall  System_GetCurrentProcessId
lcall  Pipe_Power
mov     b,a
GETX
orl    a,b
PUTX

```

Init_Done:

```

clr    a
ret

```

```

;*****
;*
;* Function Name: Pipe_Read
;*
;* Description: Function to copy data from the circular buffer
;*              back to the I/O subsystem.
;*
;* Input(s): None
;
;* Outputs(s): dpstr - pointer to input buffer
;*
;* Notes: First 4 bytes of read buffer are:
;*         length, 0, 0, 0
;*
;*****

```

```

;
; The function copies data out of the circular buffer
; and into the read return buffer. I/O allows a maximum
; return value of 250 bytes per read, so that is our upper

```

```

; bound
;

Pipe_Read:

;
; First, save *everything* away.
;
PUSH_DPTR2
PUSH_BANK_0
PUSH_BANK_1
PUSH b
PUSH acc
PUSH DPS

;
; Get the state block
;
PIPE_GET_BUFFER
;
; Save the pointer away
;
mov     R0_B1,dpl
mov     R1_B1,dph
mov     R2_B1,dpx
PUSH_DPTR1

;
; Move to the read buffer
;
mov     a,#LOW(PIPE_R_LEN)
mov     b,#HIGH(PIPE_R_LEN)
lcall  add_dpstr1_16
;
; Store the read buffer away in DPL1:DPH1:DPX1
;
mov     R3_B1,dpl
mov     R4_B1,dph
mov     R5_B1,dpx

mov     dpl1,dpl
mov     dph1,dph
mov     dpx1,dpx
POP_DPTR1
;
; Now DPL:DPH:DPX is pointing to the pipe buffer
; and DPL1:DPH1:DPX1 is pointing to the read buffer.
; Get the state vars from the pipe
;

mov     dps,#0                ; Use DPL:DPH:DPX
inc     dpstr                 ; Move to the size byte
GETX
mov     R0,a                  ; Get the start offset
inc     dpstr
GETX
mov     R1,a
inc     dpstr

;
; R7 is R0 with the high bit cleared. It will

```

```

; serve as our loop counter and as our number
; of bytes to read
;
mov     R7,R0_B0
clr     c
mov     a,R0
cjne   a,#250,$+3
jc     pipe_read_init_ret_buf
mov     R7,#250

```

```
pipe_read_init_ret_buf:
```

```

;
; We're going to need this buffer later
;
PUSH_DPTR1

;
; Ok, now the size and offset are in R0:R1.
; and DPL:DPH:DPX are pointing to the first byte
; of the pipe. Lets initialize the return buffer,
; shall we?
;
; Start by switching to DPL1:DPH1:DPX1
;
inc     dps
;
; Write the length of the return size
;
mov     a,R7
PUTX
inc     dptr
;
; Fill in the rest of the header
;
clr     a
PUTX
inc     dptr
PUTX
inc     dptr
PUTX
inc     dptr

;
; We're going to copy what we can from the pipe
; into the read buffer.
;

POP_DPTR1

;
; Looks like we're committed. Lets start reading
;
mov     dps,#1

```

```
pipe_read_test_zero:
```

```

;
; Test if we are sending zero bytes
;
mov     a,R7
jnz    pipe_read_loop

```

```
ljmp    pipe_read_exit
```

```
pipe_read_loop:
```

```
    ;
    ; Use DPL:DPH:DPX
    ;
    inc    dps
    ;
    ; Use movc for the copy
    ;
    ; Yes, I *know* how very illegal this should be
    ; as it is only possible when code space = data space,
    ; but I saw Don "The Godfather" Loomis ("Let me write you
    ; a firmware you can't refuse") do it in his code, so
    ; I claim I can to.
    ;
    mov    a,R1
    movc   a,@a+dptr

    ;
    ; Write acc into the read buffer
    ;
    inc    dps
    PUTX
    inc    dptr

    ;
    ; increment the buffer offset, decrement
    ; the size count, and loop
    ;
    inc    R1
    dec    R0
    djnz   R7,pipe_read_loop
```

```
pipe_read_exit:
```

```
    ;
    ; Get the read buffer
    ;
    mov    dpl, R0_B1
    mov    dph, R1_B1
    mov    dpX, R2_B1

    ;
    ; Update 'start' and 'offset'
    mov    dps,#0
    inc    dptr
    mov    a,R0
    PUTX
    inc    dptr
    mov    a,R1
    PUTX

    mov    dpl, R3_B1
    mov    dph, R4_B1
    mov    dpX, R5_B1
```

```

;
; Exit without incident
;

POP DPS
POP acc
POP b
POP_BANK_1
POP_BANK_0
POP_DPTR2

clr    a
ret

;*****
;*
;* Function Name: Pipe_Write
;*
;* Description:
;*
;*   Input(s): dptr  -> pointer to data to send
;*             R2    -> App Id
;*             R7    -> Thread Id
;*             R5:R4 -> length of data to send.
;*
;*   Outputs(s): a - 0 on success, exception number on failure
;*
;*
;*****

;
; This function "initiates" the write operation.
; It moves the state it has been passed into the ESB,
; and then installs a poll routine to perform the copy
; operation. The function then exits.
;

Pipe_Write:
;
; Get the state block, but make sure not
; to disrupt any of the registers we've
; been passed!
;
PUSH acc
PUSH b
PUSH dps
PUSH_BANK_0
PUSH_DPTR1

push    r4_b0
push    r5_b0
push    r2_b0
push    r7_b0
PIPE_GET_BUFFER

pop     R7_b0
pop     R2_B0
pop     R5_B0
pop     R4_B0

```

```

mov     dps,#0

;
; Write all of our state while making sure
; the scheduler doesn't run (preventing
; our poll from writing in the middle of
; our write).
;

TINIOS_ENTER_CRITICAL_SECTION

;
; Inform our poll routine that a write is occurring
;
GETX
setb    PIPE_WRITE_BIT
PUTX

;
; Fill in the state vars
;
mov     a, #LOW(PIPE_W_STATE)
mov     b, #HIGH(PIPE_W_STATE)
lcall  add_dpstr1_16

mov     a,R4
PUTX
inc     dpstr
mov     a,R5
PUTX
inc     dpstr
mov     a,R7
PUTX
inc     dpstr
mov     a,R2
PUTX

POP_DPSTR2

inc     dpstr
mov     a,dpl1
PUTX
inc     dpstr
mov     a,dph1
PUTX
inc     dpstr
mov     a,dpx1
PUTX
clr     a

mov     dpstr,#Pipe_Poll
lcall  System_RegisterPoll

TINIOS_EXIT_CRITICAL_SECTION

POP_BANK_0
pop dps

```

```

    pop b
    pop acc
    ret

;*****
;*
;* Function Name: Pipe_Poll
;*
;* Description: Poll routine installed to copy from the write
;*             buffer into the circular buffer.
;*
;* Input(s): None
;*
;* Outputs(s): a - 0 on success, exception number on failure
;*
;*
;*****

;
; This function performs the write operation. If there is not
; enough space in the circular buffer for the data, it will
; copy data into the space available, and then return. This
; will repeat until all the data has been written, where
; the poll routine calls IO_WriteFree and removes itself
; from IO_POLL
;

Pipe_Poll:
    PUSH_DPTR1
    PUSH_DPTR2
    PUSH_BANK_0
    PUSH_BANK_1
    PUSH b
    PUSH acc
    PUSH DPS

; Register Usage
; -----
; R0_B0 - Buffer size
; R1_B0 - Offset into Queue
; R2_B0 - Handle for the io subsystem
; R3_B0 - Count register for loop
; R4:R5 - Size for write (Bank 0)
; R7_B0 - Thread ID of write
; R0:R1:R2 (Bank 1) - Pointer to circular buffer offset
; R3:R4:R5 (Bank 1) - Pipe Data Structure

;
; If the write bit is not set, then
; exit out
;
    mov     dps,#0

    PIPE_GET_BUFFER
    GETX
    jb     PIPE_WRITE_BIT, pipe_poll_init
    clr    a
    ljmp   pipe_poll_exit

pipe_poll_init:

```



```

;
; Cache the pointer to the state block
;
mov     R3_B1,dpl
mov     R4_B1,dph
mov     R5_B1,dpx

mov     a,#LOW(PIPE_W_STATE)
mov     b,#HIGH(PIPE_W_STATE)
lcall  add_dpstr1_16
GETX
mov     R4,a
inc     dpstr
GETX
mov     R5,a
inc     dpstr
GETX
mov     R7,a
inc     dpstr
GETX
mov     R2,a
inc     dpstr
GETX
push    acc
inc     dpstr
GETX
push    acc
inc     dpstr
GETX
mov     dpx,a
pop     dph
pop     dpl

;
; Hold onto the write buffer
;
mov     dpl1,dpl
mov     dph1,dph
mov     dpx1,dpx

;
; Cache the pointer to the state block
;
mov     dpl,R3_B1
mov     dph,R4_B1
mov     dpx,R5_B1

;
; Get the size and the offset
;
inc     dpstr
GETX
mov     R0,a
inc     dpstr
GETX
mov     R1,a
inc     dpstr
;
; Cache the pointer to the buffer
```

```

;
mov     R0_B1,dpl
mov     R1_B1,dph
mov     R2_B1,dpx
;
; Calculate the available size for the circular
; buffer (255-size = compliment of R0) and stash
; in R3
;
mov     a,R0
cpl    a
mov     R3,a

;
; Figure out how many bytes will
; be copied in. My logic is:
; 1) If R5 is set, then size > 255. Copy in
;    what size is available
; 2) If R5 is zero and R4 > available size,
;    then copy in available size bytes
; 3) Otherwise, copy in R4 number of bytes
;
pipe_poll_calc_available:
mov     a,R5
jnz    pipe_poll_use_available

clr     c
mov     a,R3
cjne   a,R4_B0,$+3
jc     pipe_poll_use_available

mov     R3,R4_B0

pipe_poll_use_available:
;
; Check if R3 is zero, and abort if it is
;
mov     a,R3
jnz    pipe_poll_not_zero
sjmp   pipe_poll_notdone

pipe_poll_not_zero:
;
; Since we aren't going to overwrite the start
; variable, lets use it as our offset var into
; the circular buffer. Add the size to the offset
; to calculate start position
;
mov     a,R0
add     a,R1
mov     R1,a

;
; Subtract the size off R4:R5
;
clr     c
mov     a,R4

```

```

subb    a,R3
mov     R4,a
mov     a,R5
subb    a,#0
mov     R5,a

```

```
pipe_poll_loop:
```

```

;
; Yes, I *know* this is a slow approach, but
; it means I don't need to do any modulo work.
; Get a byte from the write buffer
;
inc     dps
GETX
mov     b,a
inc     dptr

;
; Return to the circular buffer
;
inc     dps

;
; Reset it, and go in offset ammount, and write
; the byte
;
mov     dpl,R0_B1
mov     dph,R1_B1
mov     dpX,R2_B1
mov     a,R1
lcall  add_dpTr1
mov     a,b
PUTX

;
; Increment counters and loop
;
inc     R1
inc     R0
djnz   R3,pipe_poll_loop

;
; If there are not any other
; bytes to write then clear the write
;
mov     a,R4
orl     a,R5
jnz     pipe_poll_notdone

mov     dpl,R3_B1
mov     dph,R4_B1
mov     dpX,R5_B1
GETX
clr     PIPE_WRITE_BIT
PUTX

lcall  System_IO_WriteFree

```

```

mov    dptr,#Pipe_Poll
lcall  System_RemovePoll

```

```

pipe_poll_notdone:

```

```

;
; Write the current state of the queue
; back into the state block
;

```

```

mov    dpl,R3_B1
mov    dph,R4_B1
mov    dpx,R5_B1
inc    dptr
mov    a,R0
PUTX

```

```

mov    a,#LOW(PIPE_W_STATE-PIPE_SIZE)
mov    b,#HIGH(PIPE_W_STATE-PIPE_SIZE)
lcall  add_dptrl_16

```

```

mov    a,R4
PUTX
inc    dptr
mov    a,R5
PUTX
inc    dptr
mov    a,R7
PUTX
inc    dptr
mov    a,R2
PUTX
inc    dptr
mov    a,dpl1
PUTX
inc    dptr
mov    a,dph1
PUTX
inc    dptr
mov    a,dpx1
PUTX

```

```

clr    a

```

```

pipe_poll_exit:

```

```

POP DPS
POP acc
POP b
POP_BANK_1
POP_BANK_0
POP_DPTR2
POP_DPTR1

```

```

clr    a
ret

```

```

;*****
;*

```

```

;* Function Name: Pipe_Open
;*
;* Description: Function to initialize the driver
;*
;* Input(s): a      -> Port Number
;*           R2     -> Driver Handle
;*
;* Outputs(s): a - 0 if success, exception number otherwise
;*
;* Notes: I am using the Native Library init routine to
;*        create the state block, so we will just initialize
;*        it in this function.
;*
;*****

```

Pipe_Open:

```

;
; Initialize the pipe data structure
;
PIPE_GET_BUFFER
clr    a
PUTX
inc    dptr
PUTX
inc    dptr
PUTX

;
; Clear out anything that remains in the pipe read buffer
;
mov    a,#LOW(PIPE_R_LEN-PIPE_START)
mov    b,#HIGH(PIPE_R_LEN-PIPE_START)
lcall  add_dpstr1_16
clr    a
PUTX
inc    dptr
PUTX
inc    dptr
PUTX
inc    dptr
PUTX
inc    dptr

;
; Increment the reference count
;

clr    a
ret

;*****
;*
;* Function Name: Pipe_Close
;*
;* Description: Function to shutdown the driver
;*
;* Input(s): a      -> Port Number
;*
;* Outputs(s): a - 0 if success, exception number otherwise
;*
;* Notes: Simply making sure the Poll routine has been removed

```

```

;*          before exiting.
;*
;*
;*****

```

Pipe_Close:

```

;
; First, remove the poll routine if it is resident
;
mov     dptr,#Pipe_Poll
lcall  System_RemovePoll

```

```

clr     a
ret

```

```

;*****
;*
;* Function Name: Pipe_Ioctl
;*
;* Description: Use to make driver-specific calls
;*
;* Input(s):  dptr  -> pointer to argument array
;*           R2    -> Driver Handle
;*           R5:R4 -> length of data to send.
;*
;* Outputs(s): a - 0 if success, exception number otherwise
;*
;* Notes:
;*
;*****

```

Pipe_IOCTL:

```

clr     a
ret

```

```

;*****
;*
;* Function Name: Pipe_Available
;*
;* Description: Use to make driver-specific calls
;*
;* Input(s):  R2    -> Driver Handle
;*
;* Outputs(s): a - 0 if success, exception number otherwise
;*           R0:R1:R2:R2 - Size of available pipe data.
;*
;* Notes:
;*
;*****

```

```

;
; Pulls the size out of the ESB and returns it
;

```

Pipe_Available:

```

PUSH_DPTR1

```

 PUSH_BANK_0

PUSH acc

PUSH b

PIPE_GET_BUFFER

TINIOS_ENTER_CRITICAL_SECTION

inc dptr

GETX

TINIOS_EXIT_CRITICAL_SECTION

mov R0_B1,a

clr a

mov R1_B1,a

mov R2_B1,a

mov R3_B1,a

POP b

POP acc

POP_BANK_0

POP_DPTR1

ret

```

;*****
;*
;* Function Name: Pipe_Init
;*
;* Description: Prepares the driver for registration
;*
;* Input(s): None
;*
;* Outputs(s): All registration registers are set.
;*
;*
;*****

```

Pipe_Init:

mov dptr,#Pipe_Read

mov R0_B0,dpl

mov R1_B0,dph

mov R0_B1,dpx

mov dptr,#Pipe_Write

mov R2_B0,dpl

mov R3_B0,dph

mov R2_B1,dpx

mov dptr,#Pipe_Open

mov R4_B0,dpl

mov R5_B0,dph

mov R4_B1,dpx

mov dptr,#Pipe_Close

mov R6_B0,dpl

mov R7_B0,dph

mov R6_B1,dpx

mov dptr,#Pipe_IOCTL

```

    mov R0_B2,dpl
    mov R1_B2,dph
    mov R0_B3,dpx

    mov dptr,#Pipe_Available
    mov R2_B2,dpl
    mov R3_B2,dph
    mov R2_B3,dpx
    ret

;*****
;*
;* Function Name: Pipe_Power
;*
;* Description: Helper method - performs a 2^a operation
;*
;* Input(s): acc - value
;*
;* Outputs(s): acc - Two to the power of parameter (8 bits only)
;*
;*
;*****

Pipe_Power:
    push    b
    mov     b,a
    mov     a,#1
    jz      pipe_power_exit
pipe_power_loop:
    rl     a
    djnz   b,pipe_power_loop

pipe_power_exit:
    pop     b
    ret

;*****
;*
;* Function Name: Pipe_ProcessExit
;*
;* Description: Called on process destruction. Cleans up
;*              things left behind.
;*
;* Input(s): acc - Process ID
;*
;* Outputs(s): none
;*
;* Note(s): PID is passed in R0
;*
;*****

Pipe_ProcessExit:
    PUSH_DPTR1
    PUSH_BANK_0
    push    b
    push    R0_B1
    push    R1_B1
    push    R2_B1

    PIPE_GET_BUFFER

```



```

jz      pipe_processexit_start
ljmp    pipe_processexit_exit

```

```
pipe_processexit_start:
```

```

;
; Hold onto the pointer
;
mov     R0_B1,dpl
mov     R1_B1,dph
mov     R2_B1,dpx
;
; Hold onto the handle
;
mov     R6,R2_B0
mov     R7,R3_B0
;
; Hold the PID
;
;
; We need to clear the process bit
;
mov     a,#LOW(PIPE_REFCOUNT)
mov     b,#HIGH(PIPE_REFCOUNT)
lcall   add_dpstr1_16
;
; Set the bit denoting the process id
;
mov     a,R0
lcall   Pipe_Power

cpl     a
mov     b,a
GETX
anl     a,b
PUTX

;
; If no one else is holding onto this driver,
; then perform the proper cleanup.
;
jz      pipe_processexit_cleanup
sjmp    pipe_processexit_exit

```

```
pipe_processexit_cleanup:
```

```

;
; Since we're still in the neighborhood,
; lets get the driver num and store it away
;
inc     dpstr
GETX
;
; First, remove the driver from the driver table
;
lcall   System_Unregister_Driver

;
; Next, remove the poll routine if it is resident
;
mov     dpstr,#Pipe_Poll

```

```

    lcall    System_RemovePoll

    ;
    ; Remove ourselves from the process destroy
    ; callback
    ;
    mov     dptr, #Pipe_ProcessExit
    lcall   System_UnregisterProcessDestroyFunction

    ;
    ; Finally, lets clear the ESB
    ;
    mov     dptr,#QueueID
    lcall   NatLib_RemoveEphemeralStateBlock

    mov     R2,R6_B0
    mov     R3,R7_B0
    lcall   MM_Free
pipe_processexit_exit:
    pop     R2_B1
    pop     R1_B1
    pop     R0_B1
    pop     b
    POP_BANK_0
    POP_DPTR1
    clr     a
    ret

;
; TINIconvertor in version 1.02 requires a native method
;
Native_GetPipeDriver:
    PIPE_GET_BUFFER
    mov     a,#LOW(PIPE_DRVNUM)
    mov     b,#HIGH(PIPE_DRVNUM)
    lcall   add_dpstr1_16
    GETX
    mov     R0,a
    clr     a
    mov     R1,a
    mov     R2,a
    mov     R3,a
    ret

QueueID:
db "DSPIPEEX",0

Debug_Str:
db "Debug Point ",0

```

END

Appendix B: driver.inc

```

#include(tinidriver.inc)

;
;
; Structure for my buffer
;
; struct PipeQueue
; {
;     //
;     // Circular buffer state vars
;     //
;     ul  flags
;     ul  size
;     ul  start
;     ul  buffer[256]
;     //
;     // Read return value
;     //
;     ul  r_len
;     ul  r_reserved[3]
;     ul  r_buffer[256]
;     //
;     // Write State variables
;     //
;     ul  w_state[8]
;     //
;     // Keep track of how many times we have been opened
;     //
;     ul  reference_count
;     //
;     // The driver number we loaded into
;     //
;     ul  driver_num
; }

PIPE_FLAGS      equ 0
PIPE_SIZE       equ PIPE_FLAGS+1
PIPE_START      equ PIPE_SIZE+1
PIPE_BUFFER     equ PIPE_START+1
PIPE_R_LEN      equ PIPE_BUFFER+256
PIPE_R_RESV     equ PIPE_R_LEN+1
PIPE_R_BUFFER   equ PIPE_R_LEN+3
PIPE_W_STATE    equ PIPE_R_BUFFER+256
PIPE_REFCOUNT   equ PIPE_W_STATE+8
PIPE_DRVNUM     equ PIPE_REFCOUNT+1
PIPE_END        equ PIPE_DRVNUM+1

;
; Used for denoting a write in progress
; in the flags byte
;
PIPE_WRITE_BIT  equ      acc.7

;
; Get the pipe buffer from TINIOS

```

```
;
PIPE_GET_BUFFER MACRO
    mov    dptr, #QueueID          ; see if we have an ESB installed
    lcall NatLib_GetEphemeralStateBlock
ENDM
```

```
;
; Debug Macro
;
PIPE_DEBUG MACRO PARAM value
    PUSH_DPTR1
    push  acc
    push  dps
    mov   dps, #0
    mov   dptr, #Debug_Str
    lcall info_sendstring
    mov   a, #value
    lcall info_sendtwohex
    lcall info_sendcrlf
    pop   dps
    pop   acc
    POP_DPTR1
ENDM
```

Appendix C: PipeDriver.java

```
package com.dalsemi.pipeio;

import com.dalsemi.comm.*;

public class PipeDriver
{
    static int pipe_port;
    //
    // This method dynamically looks up the
    // next available driver from the system
    // and allocates it, or uses the current
    // driver number if the pipe is
    // already loaded.
    //
    static native int getPipeDriver();

    static
    {
        System.loadLibrary("driver.tlib");
        pipe_port = getPipeDriver();
    }
}
```

Appendix D: PipeInputStream.java

```
package com.dalsemi.pipeio;

import com.dalsemi.comm.*;
import java.io.*;

public class PipeInputStream extends InputStream
{
    //
    // Handle to the pipe input stream
    //
    int pipe_handle;
    //
    // Input buffer for single byte reads
    //
    byte [] charArr;

    public PipeInputStream() throws Exception
    {
        //
        // Denote the pipe as invalid
        //
        pipe_handle = -1;
        charArr = new byte[1];
        try
        {
            //
            // On any exception, make the pipe invalid
            //
            pipe_handle = NativeComm.open(PipeDriver.pipe_port,
                NativeComm.STREAM_STDIN);
        }
        catch (Exception E)
        {
            pipe_handle = -1;
            throw E;
        }
    }

    public int read(byte arr[], int offset, int length) throws IOException
    {
        //
        // Perform a validity check
        //
        if (pipe_handle == -1)
            throw new IOException("pipe is not valid");
        //
        // Perform a length/offset check. This is *not* done
        // by I/O for you.
        //
        if ((offset >= 0) && ((offset+length) <= arr.length))
        {
            return NativeComm.read(
                pipe_handle, arr, offset, length, 0, false);
        }
        else throw new ArrayIndexOutOfBoundsException(
            "offset = " + offset + "length = " + length);
    }
}
```

```
public int read() throws IOException
{
    int len;
    //
    // Perform a validity check
    //
    if (pipe_handle == -1)
        throw new IOException("pipe is not valid");
    //
    // Read a single byte
    //
    len = NativeComm.read(pipe_handle, charArr, 0, 1, 0, false);

    return ((len==1)?((int)charArr[0]&0xFF):-1);
}

public void close() throws IOException
{
    if (pipe_handle != -1)
        NativeComm.close(pipe_handle);
}
}
```

Appendix E: PipeOutputStream.java

```

package com.dalsemi.pipeio;

import com.dalsemi.comm.*;
import java.io.*;

public class PipeOutputStream extends OutputStream
{
    //
    // Handle to the pipe input stream
    //
    int pipe_handle;
    //
    // Array for passing single byte values
    //
    byte [] charArr;

    public PipeOutputStream() throws Exception
    {
        //
        // Denote the pipe as invalid
        //
        pipe_handle = -1;
        charArr = new byte[1];
        try
        {
            //
            // On any exception, make the pipe invalid
            //
            pipe_handle = NativeComm.open(PipeDriver.pipe_port,
                NativeComm.STREAM_STDOUT);
        }
        catch (Exception E)
        {
            pipe_handle = -1;
            throw E;
        }
    }

    public void write(byte [] arr, int offset, int len) throws IOException
    {
        //
        // Perform a validity check
        //
        if (pipe_handle == -1)
            throw new IOException("pipe is not valid");
        //
        // Perform a length/offset check. This is *not* done
        // by I/O for you.
        //
        if ((offset >= 0) && ((offset+len) <= arr.length))
        {
            NativeComm.write(pipe_handle, arr, offset, len);
        }
        else
            throw new ArrayIndexOutOfBoundsException(
                "offset = " + offset + " len = " + len);
    }

    public void write(int x) throws IOException

```

```
{
    //
    // Perform a validity check
    //
    if (pipe_handle == -1)
        throw new IOException("pipe is not valid");

    charArr[0] = (byte)(x & 0xFF);

    NativeComm.write(pipe_handle, charArr, 0, 1);
}

public void close() throws IOException
{
    if (pipe_handle != -1)
        NativeComm.close(pipe_handle);
}
}
```